PETSc Tutorial

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with thanks to Jed Brown, Matt Knepley, Karl Rupp, and Barry Smith for slides additional tutorial material available via https://www.mcs.anl.gov/petsc

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Outline

PETSc Tutorial

Philosophy Vectors and matrices (Vec, Mat) Linear solvers (KSP, PC) Nonlinear solvers (SNES) DAE/ODE integrators: Timestepping (TS) Optimization solvers (TAO) Topology abstractions: Distributed arrays (DMDA) Understanding performance



About PETSc

PETSc Origins

PETSc was developed as a Platform for Experimentation

We want to experiment with different

Models

Discretizations

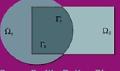
Solvers

Algorithms

These boundaries are often blurred...



Parallel Multilevel Methods for Elliptic Partial Differential Equations,



Barry Smith, Petter Bjørstad, and William Gropp

Portable Extensible Toolkit for Scientific Computing

Architecture

Tightly coupled (e.g. XT5, BG/P, Earth Simulator) Loosely coupled such as network of workstations GPU clusters (many vector and sparse matrix kernels)

Software Environment

Operating systems (Linux, Mac, Windows, BSD, proprietary Unix) Any compiler

Usable from C, C++, Fortran 77/90, Python

Real/complex, single/double/quad precision, 32/64-bit int

System Size

500B unknowns, 75% weak scalability on Jaguar (225k cores) and Jugene (295k cores)

Same code runs performantly on a laptop

Free to everyone (BSD license), open development

Portable Extensible Toolkit for Scientific Computing

Philosophy: Everything has a plugin architecture

Vectors, Matrices, Coloring/ordering/partitioning algorithms Preconditioners, Krylov accelerators Nonlinear solvers, Time integrators Spatial discretizations/topology

Example

Vendor supplies matrix format and associated preconditioner, distributes compiled shared library.

Application user loads plugin at runtime, no source code in sight.

Portable Extensible Toolkit for Scientific Computing

Toolset

Algorithms (Parallel) debugging aids Low-overhead profiling

Composability

Try new algorithms by choosing from product space Composing existing algorithms (multilevel, domain decomposition, splitting)

Experimentation

Impossible to pick the solver *a priori* PETSc's response: expose an algebra of composition Keep solvers decoupled from physics and discretization

Portable Extensible Toolkit for Scientific Computing

Computational Scientists

PyLith (CIG), Underworld (Monash), Magma Dynamics (LDEO, Columbia), PFLOTRAN (DOE), SHARP/UNIC (DOE)

Algorithm Developers (iterative methods and preconditioning)

Package Developers

SLEPc, TAO, Deal.II, Libmesh, FEniCS, PETSc-FEM, MagPar, OOFEM, FreeCFD, OpenFVM

Funding

Department of Energy SciDAC, ASCR ISICLES, MICS Program, INL Reactor Program National Science Foundation CIG, CISE, Multidisciplinary Challenge Program

Documentation and Support

Hundreds of tutorial-style examples

Hyperlinked manual, examples, and manual pages for all routines

Support at petsc-maint@mcs.anl.gov petsc-users@mcs.anl.gov

Developing parallel, nontrivial PDE solvers that deliver high performance is still difficult and requires months (or even years) of concentrated effort.

PETSc is a toolkit that can ease these difficulties and reduce the development time, but it is not a black-box PDE solver, nor a silver bullet.

- Barry Smith

Obtaining PETSc

Linux Package Managers Web: http://mcs.anl.gov/petsc, download tarball Git: https://bitbucket.org/petsc/petsc

Installing PETSc

```
$> cd /path/to/petsc/workdir
$> git clone https://bitbucket.org/petsc/petsc -b maint
$> cd petsc
```

```
$> export PETSC_DIR=$PWD PETSC_ARCH=mpich-gcc-dbg
$> ./configure
    --with-cc=gcc --with-fc=gfortran --with-cxx=g++
    --download-fblaslapack --download-mpich
    --download-{ml,hypre,superlu}
```

Most packages can be automatically

Downloaded

Configured and Built (in \$PETSC_DIR/\$PETSC_ARCH/externalpackages)
Installed with PETSC

Works for (list incomplete)

petsc4py

PETSc documentation utilities (Sowing, Igrind, c2html)

BLAS, LAPACK, BLACS, ScaLAPACK

MPICH, MPE, OpenMPI

ParMetis, Chaco, Party, Scotch, Zoltan

MUMPS, Spooles, SuperLU, SuperLU_Dist, UMFPack, pARMS

PaStiX, BLOPEX, FFTW, SPRNG

HYPRE, ML, SPAI

SUNDIALS

Triangle, TetGen, FIAT

HDF5, Boost

Commits to the PETSc Repository

Graph of commit histroy



Sep 18, 1994 - Aug 4, 2016

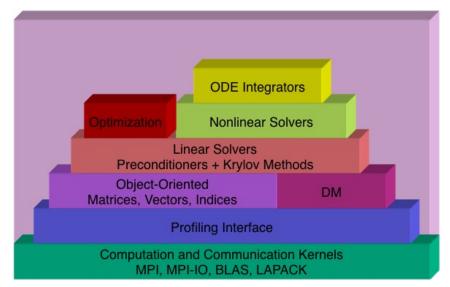
Contributions: Commits -

Contributions to master, excluding merge commits



PETSc Pyramid

PETSc Structure



Library Design

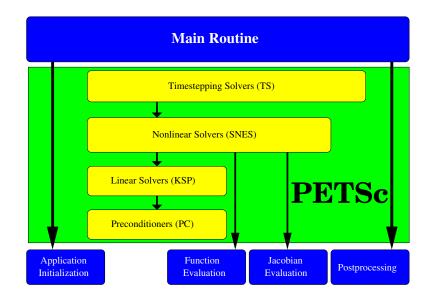
Numerical libraries should interact at a higher level than MPI

MPI coordinates data movement and synchronization for data parallel applications

Numerical libraries should coordinate access to a given data structure MPI can handle data parallelism and something else (runtime engine) handle task parallelism

Algorithm should be data structure neutral, but its main operation is still to structure access

Flow Control for a PETSc Application



Initialization

Call PetscInitialize()

Setup static data and services Setup MPI if it is not already Can set PETSC_COMM_WORLD to use your communicator (can always use subcommunicators for each object)

Call PetscFinalize()

Calculates logging summary Can check for memory leaks, unused options Shutdown and release resources

Recommend initializing PETSc only once

Every object in PETSc supports a basic interface

Function	Operation
Create()	create the object
Get/SetName()	name the object
Get/SetType()	set the implementation type
Get/SetOptionsPrefix()	set the prefix for all options
SetFromOptions()	customize object from the command line
SetUp()	preform other initialization
View()	view the object
Destroy()	cleanup object allocation

Also, all objects support the -help option.



Vectors and Matrices

Vector Algebra

What are PETSc vectors?

Fundamental objects representing field solutions, right-hand sides, etc. Each process locally owns a subvector of contiguous global data

How do I create vectors?

VecCreate(MPI_Comm, Vec *)
VecSetSizes(Vec, int n, int N)
VecSetType(Vec, VecType typeName)
VecSetFromOptions(Vec) - Can set the type at runtime

A PETSc Vec

Has a direct interface to the values
Supports all vector space operations
 VecDot(), VecNorm(), VecScale()
Has unusual operations, e.g. VecSqrt(), VecWhichBetween()
Communicates automatically during assembly
Has customizable communication (scatters)

Collectivity

MPI communicators (MPI_Comm) specify collectivity Processes involved in a computation Constructors are collective over a communicator

VecCreate(MPI_Comm comm, Vec *x)
Use PETSC_COMM_WORLD for all processes and PETSC_COMM_SELF for one

Some operations are collective, while others are not

collective: VecNorm()
not collective: VecGetLocalSize()

Sequences of collective calls must be in the same order on each process

Parallel Assembly

Vectors and Matrices

Processes may set an arbitrary entry Must use proper interface Entries need not be generated locally Local meaning the process on which they are stored PETSc automatically moves data if necessary Happens during the assembly phase

Vector Assembly

A three step process

Each process sets or adds values Begin communication to send values to the correct process Complete the communication

```
VecSetValues(Vec v, int n, int rows[], PetscScalar
values[], mode)
```

mode is either INSERT_VALUES or ADD_VALUES

Two phase assembly allows overlap of communication and computation

```
VecAssemblyBegin(Vec v)
VecAssemblyEnd(Vec v)
```

```
VecGetSize(x, &N);
MPI_Comm_rank(PETSC_COMM_WORLD, &rank);
if (rank == 0) {
  for(i = 0, val = 0.0; i < N; i++, val += 10.0) {
    VecSetValues(x, 1, &i, &val, INSERT_VALUES);
  }
}
/* These routines ensure that the data is distributed to the
other processes */
VecAssemblyBegin(x);
VecAssemblyEnd(x);
```

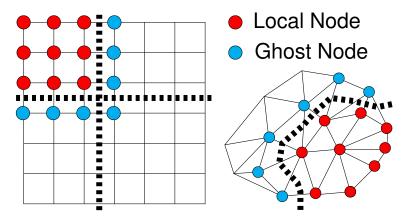
```
VecGetOwnershipRange(x, &low, &high);
for(i = low,val = low*10.0; i < high; i++,val += 10.0) {
    VecSetValues(x, 1, &i, &val, INSERT_VALUES);
}
/* These routines ensure that the data is distributed to the
other processes */
VecAssemblyBegin(x);
VecAssemblyEnd(x);
```

Ghost Values

To evaluate a local function f(x), each process requires

its local portion of the vector x

its ghost values, bordering portions of *x* owned by neighboring processes



It is sometimes more efficient to directly access local storage of a Vec. PETSc allows you to access the local storage with VecGetArray (Vec, double *[]) You must return the array to PETSc when you finish VecRestoreArray (Vec, double *[]) Allows PETSc to handle data structure conversions Commonly, these routines are inexpensive and do not involve a copy

Matrices

Definition (Matrix)

A matrix is a linear transformation between finite dimensional vector spaces.

Definition (Forming a matrix)

Forming or assembling a matrix means defining its action in terms of entries (usually stored in a sparse format).

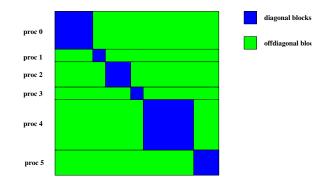
```
MatCreate(MPI_Comm, Mat *)
MatSetSizes(Mat, int m, int n, int M, int N)
MatSetType (Mat, MatType typeName)
MatSetFromOptions (Mat)
    Can set the type at runtime
MatMPIBAIJSetPreallocation(Mat,...)
    important for assembly performance
MatSetBlockSize(Mat, int bs)
    for vector problems
MatSetValues (Mat, ...)
    MUST be used, but does automatic communication
   MatSetValuesLocal, MatSetValuesStencil,
   MatSetValuesBlocked
```

The PETSc Mat has a single user interface, Matrix assembly MatSetValues() Matrix-vector multiplication MatMult() Matrix viewing MatView() but multiple underlying implementations. AIJ, Block AIJ, Symmetric Block AIJ, Dense, Elemental Matrix-Free etc.

A matrix is defined by its interface, not by its data structure.

Parallel Sparse Matrix

Each process locally owns a submatrix of contiguous global rows Each submatrix consists of diagonal and off-diagonal parts



MatGetOwnershipRange(Mat A, int *start, int *end)
start: first locally owned row of global matrix
end-1: last locally owned row of global matrix

Simple 3-point stencil for 1D Laplacian

```
v[0] = -1.0; v[1] = 2.0; v[2] = -1.0;
if (rank == 0) {
  for(row = 0; row < N; row++) {
    cols[0] = row-1; cols[1] = row; cols[2] = row+1;
    if (row == 0) {
      MatSetValues(A, 1, &row, 2, &cols[1], &v[1],
                   INSERT VALUES):
    } else if (row == N-1) {
      MatSetValues(A,1,&row,2,cols,v,INSERT VALUES);
    } else {
      MatSetValues(A,1,&row,3,cols,v,INSERT VALUES);
MatAssemblyBegin(A, MAT FINAL ASSEMBLY);
MatAssemblyEnd(A, MAT FINAL ASSEMBLY);
```

A More Efficient Way

Advantages

All ranks busy: Scalable! Amount of code essentially unchanged



Iterative Solvers

Matrices

What can we do with a matrix that doesn't have entries?

Krylov solvers for Ax = b

Krylov subspace: $\{b, Ab, A^2b, A^3b, \dots\}$

Convergence rate depends on the spectral properties of the matrix

For any popular Krylov method \mathcal{K} , there is a matrix of size *m*, such that \mathcal{K} outperforms all other methods by a factor at least $\mathcal{O}(\sqrt{m})$ [Nachtigal et. al., 1992]

Typically...

The action $y \leftarrow Ax$ can be computed in $\mathcal{O}(m)$

Aside from matrix multiply, the n^{th} iteration requires at most $\mathcal{O}(mn)$

PETSc Solvers

Linear Solvers - Krylov Methods

Using PETSc linear algebra, just add:

KSPSetOperators(KSP ksp, Mat A, Mat M)
KSPSolve(KSP ksp, Vec b, Vec x)

Can access subobjects

KSPGetPC(KSP ksp, PC *pc)

Preconditioners must obey PETSc interface Basically just the KSP interface

Can change solver dynamically from the command line, -ksp_type

Linear Solvers in PETSc KSP

Linear Solvers in PETSc KSP (Excerpt)

Richardson

Chebychev

Conjugate Gradient

BiConjugate Gradient

Generalized Minimum Residual Variants

Transpose-Free Quasi-Minimum Residual

Least Squares Method

Conjugate Residual

Complete table of solvers



Preconditioners

Preconditioning

Idea: improve the conditioning of the Krylov operator

Left preconditioning

$$(P^{-1}A)x = P^{-1}b$$

{ $P^{-1}b, (P^{-1}A)P^{-1}b, (P^{-1}A)^2P^{-1}b, \dots$ }

Right preconditioning

$$(AP^{-1})Px = b$$

 $[b, (P^{-1}A)b, (P^{-1}A)^2b, \dots]$

The product $P^{-1}A$ or AP^{-1} is *not* formed.

A preconditioner \mathcal{P} is a method for constructing a matrix (just a linear function, not assembled!) $P^{-1} = \mathcal{P}(A, A_p)$ using a matrix A and extra information A_p , such that the spectrum of $P^{-1}A$ (or AP^{-1}) is well-behaved.

Preconditioners in PETSC PC

Preconditioners in PETSc PC (Excerpt)

Jacobi block Jacobi SOR Additive Schwarz Incomplete factorizations (ILU(k), ICC(k)) Multigrid (geometric, algebraic) Physics-based splitting Approximate inverses Substructuring Matrix-free (infrastructure for custom approaches provided by user)

Complete table of solvers

Splitting for Multiphysics

$$\begin{bmatrix} A & B \\ C & D \end{bmatrix} \begin{bmatrix} x \\ y \end{bmatrix} = \begin{bmatrix} f \\ g \end{bmatrix}$$

Relaxation: -pc_fieldsplit_type [additive,multiplicative,symmetric_multiplicative]

$$\begin{bmatrix} A \\ & D \end{bmatrix}^{-1} \begin{bmatrix} A \\ C \end{bmatrix}^{-1} \begin{bmatrix} A \\ & \mathbf{I} \end{bmatrix}^{-1} \begin{bmatrix} A \\ & \mathbf{I} \end{bmatrix}^{-1} \begin{pmatrix} \mathbf{I} - \begin{bmatrix} A & B \\ & \mathbf{I} \end{bmatrix} \begin{bmatrix} A \\ C \end{bmatrix}^{-1} \begin{pmatrix} \mathbf{I} \end{bmatrix}^{-1} \begin{pmatrix} \mathbf{I} \end{bmatrix} \begin{bmatrix} A \\ & \mathbf{I} \end{bmatrix}^{-1} \begin{pmatrix} \mathbf{I} \end{bmatrix} \begin{bmatrix} A \\ & \mathbf{I} \end{bmatrix}^{-1} \begin{pmatrix} \mathbf{I} \end{bmatrix}^{-1} \begin{pmatrix} \mathbf{I} \end{bmatrix}^{-1} \begin{pmatrix} \mathbf{I} \end{bmatrix} \begin{pmatrix} \mathbf{I} \end{bmatrix}^{-1} \begin{pmatrix} \mathbf{I} \end{pmatrix}^{-1} \end{pmatrix}^{-1} \begin{pmatrix} \mathbf{I} \end{pmatrix}^{-1} \begin{pmatrix} \mathbf{I}$$

Gauss-Seidel inspired, works when fields are loosely coupled Factorization: -pc_fieldsplit_type schur

$$\begin{bmatrix} A & B \\ S \end{bmatrix}^{-1} \begin{bmatrix} \mathbf{1} \\ CA^{-1} & \mathbf{1} \end{bmatrix}^{-1}, \qquad S = D - CA^{-1}B$$

robust (exact factorization), can often drop lower block how to precondition *S* which is usually dense?

interpret as differential operators, use approximate commutators

Solver Composition

Unintrusive composition of multigrid and block preconditioning

We can build many preconditioners from the literature on the command line

User code does not depend on matrix format, preconditioning method, nonlinear solution method, time integration method (implicit or IMEX), or size of coupled system (except for driver).



Nonlinear Solvers

Nonlinear solvers in PETSc SNES

LS, TR Newton-type with line search and trust region NRichardson Nonlinear Richardson, usually preconditioned VIRS, VISS reduced space and semi-smooth methods for variational inequalities QN Quasi-Newton methods like BFGS NGMRES Nonlinear GMRES NCG Nonlinear Conjugate Gradients GS Nonlinear Gauss-Seidel/multiplicative Schwarz sweeps

- FAS Full approximation scheme (nonlinear multigrid)
- MS Multi-stage smoothers, often used with FAS for hyperbolic problems

Shell Your method, often used as a (nonlinear) preconditioner

We will illustrate basic solver usage with SNES.

Use ${\tt SNESSetFromOptions}$ () so that everything is set dynamically

Use <code>-snes_type</code> to set the type or take the default

Override the tolerances

```
Use -snes_rtol and -snes_atol
```

View the solver to make sure you have the one you expect

```
Use -snes_view
```

For debugging, monitor the residual decrease

```
Use -snes_monitor
```

Use -ksp_monitor to see the underlying linear solver

Newton iteration: Workhorse of SNES

Standard form of a nonlinear system

F(u) = 0



Solve:
$$J(u)w = -F(u)$$

Update: $u^+ \leftarrow u + w$

Quadratically convergent near a root: $|u^{n+1} - u^*| \in O(|u^n - u^*|^2)$ Picard is the same operation with a different J(u)

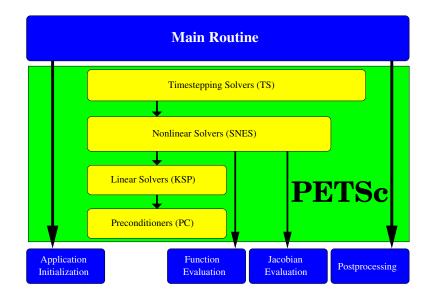
Example (Nonlinear Poisson)

$$F(u) = 0 \quad \sim \quad - \div \left[(1+u^2)\nabla u \right] - f = 0$$

$$J(u)w \quad \sim \quad - \div \left[(1+u^2)\nabla w + 2uw\nabla u \right]$$



Flow Control for a PETSc Application



SNES Paradigm

SNES Interface based upon Callback Functions

FormFunction(), set by SNESSetFunction()
FormJacobian(), set by SNESSetJacobian()

Evaluating the nonlinear residual F(x)

Solver calls the user's function

User function gets application state through the ${\tt ctx}$ variable

PETSc never sees application data

SNES Function

F(u)=0

The user provided function which calculates the nonlinear residual has signature

- $\scriptstyle \times$ The current solution
- r The residual
- ctx The user context passed to SNESSetFunction()
 - Use this to pass application information, e.g. physical constants

SNES Jacobian

User-provided function calculating the Jacobian Matrix

- $\scriptstyle \times$ The current solution
- $\ensuremath{\mathbb{J}}$ The Jacobian
- M The Jacobian preconditioning matrix (possibly J itself)
- ctx The user context passed to SNESSetFunction()
 - Use this to pass application information, e.g. physical constants
- Possible MatStructure values are:

```
SAME_NONZERO_PATTERN
DIFFERENT_NONZERO_PATTERN
```

Alternatives

a builtin sparse finite difference approximation ("coloring") automatic differentiation (ADIC/ADIFOR)



Time Integration

ODE forms supported

$$G(t, x, \dot{x}) = F(t, x)$$

$$J_{\alpha} = \alpha G_{\dot{x}} + G_{x} \text{ or }$$

$$M(t)\dot{x} = F(t, x)$$

$$J_{\alpha} = \alpha M \text{ or }$$

$$\dot{x} = F(t, x)$$

User provides:

FormRHSFunction (ts, t, x, F, void *ctx); FormIFunction (ts, t, x, \dot{x}, G , void *ctx); FormIJacobian (ts, $t, x, \dot{x}, \alpha, J, J_p$, void *ctx);

Explicit methods are easy and accurate, but must resolve all time scales

Reactions, acoustics, incompressibility

Implicit methods are robust

Mathematically good for stiff systems Harder to implement, need efficient solvers

Implicit-explicit methods blend the benefits of both

Benefits of explicit for parts of problems that are not stiff Benefits of implicit where needed Good approach for multiphysics applications

Some TS Methods

TSSSPRK104 10-stage, fourth order, low-storage, optimal explicit SSP Runge-Kutta $c_{\text{eff}} = 0.6$ (Ketcheson 2008)

TSARKIMEX2E second order, one explicit and two implicit stages, *L*-stable, optimal (Constantinescu)

TSARKIMEX3 (and 4 and 5), L-stable (Kennedy and Carpenter, 2003)

TSROSWRA3PW three stage, third order, for index-1 PDAE, *A*-stable, $R(\infty) = 0.73$, second order strongly *A*-stable embedded method (Rang and Angermann, 2005)

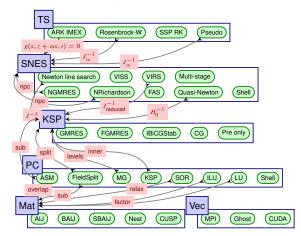
TSROSWRA34PW2 four stage, third order, *L*-stable, for index 1 PDAE, second order strongly *A*-stable embedded method (Rang and Angermann, 2005)

TSROSWLLSSP3P4S2C four stage, third order, *L*-stable implicit, SSP explicit, *L*-stable embedded method (Constantinescu)

TS solver options

Interface to highest level that makes sense for your problem: TS preferable

Customize algorithmic layers as needed: SNES, KSP, PC Can make choices at runtime for algorithms and data structures



Be willing to experiment with algorithms

No optimality without interplay between physics and algorithmics

Adopt flexible, extensible programming

Algorithms and data structures not hardwired

Be willing to play with the real code

Toy models have limited usefulness But make test cases that run quickly

If possible, profile before integration Automatic in PETSc PETSc does not seize main(), does not control output Propogates errors from underlying packages, flexible Nothing special about MPI_COMM_WORLD

Can wrap existing data structures/algorithms

MatShell, PCShell, full implementations

VecCreateMPIWithArray()

MatCreateSeqAIJWithArrays()

Use an existing semi-implicit solver as a preconditioner

Usually worthwhile to use native PETSc data structures unless you have a good reason not to

Uniform interfaces across languages

C, C++, Fortran 77/90, Python

Do not have to use high level interfaces (e.g. SNES, TS, DM) but PETSc can offer more if you do, like MFFD and SNES Test Use the package with the highest level of abstraction that uses PETSc

Eigenvalues - SLEPc

Finite Elements - Deal.II, Libmesh, FEniCS, PETSc-FEM, OOFEM,

Finite Elements and Multiphysics - MOOSE

Finite Volumes - FreeCFD, OpenFVM

Wave Propagation - PyClaw

Micromagnetics - MagPar

Numerical Optimization - Toolkit for Advanced Optimization (TAO)

TAO is now part of PETSc



Nonlinear Optimization: TAO

Developers: Todd Munson, Jason Sarich, Stefan Wild

Nonlinear Optimization

Solves Nonlinear Optimization Problems:

$$f: \mathbb{R}^N \mapsto \mathbb{R}$$
$$\min_{x \in \mathbb{R}^N} f(x)$$

With optional variable bounds:

subject to
$$x_l \leq x \leq x_u$$
 (bounds)

Or complementarity constraints:

$$F_i(x^*) \ge 0 \qquad \text{if } x_i^* = \ell_i$$

$$F_i(x^*) = 0 \qquad \text{if } \ell_i < x_i^* < u_i$$

$$F_i(x^*) \le 0 \qquad \text{if } x_i^* = u_i.$$

Also some support for PDE-constrained applications and general contraints

TAO Algorithms

TAO provides a suite of (iterative) nonlinear optimization algorithms. Typically, each iteration involves calculating a *search direction* d_k , then function values and gradients along that direction are calculated until desired conditions are met.

Newton's Method

Calculate the direction d_{k+1} by solving the system:

$$\nabla^2 f(x_k) d_{k+1} = -\nabla f(x_k)$$

Quasi-Newton Methods

Use approximate Hessian $B_k \approx \nabla^2 f(x_k)$. Choose a formula for B_k so that B_k relies on first derivative information only, can be easily stored and $B_k d_{k+1} = -\nabla f(x_k)$ can be easily solved.

Conjugate Gradient

Derivative Free

TAO Solvers

Solvers available in TAO

	handles constraints	requires gradient	requires Hessian
Quasi-Newton (Imvm)	no	yes	no
Newton Line Search (nls)	no	yes	yes
Newton Trust Region (ntr)	no	yes	yes
Newton Trust with Line Search (ntl)	no	yes	yes
Conjugate Gradient (cg)	no	yes	no
Nelder-Mead (nm)	no	no	no
Quasi-Newton (blmvm)	bounds	yes	no
Newton Trust Region (tron)	bounds	yes	yes
Conjugate Gradient (gpcg)			
(Quadratic objective only)	bounds	yes	no
Model-based derivative free			
nonlinear least-squares (pounders)	yes	no	no
Semismooth – Feasibility-enforced			
(SSFLS)	complementarity	yes	yes
Semismooth – Feasibility not enforced			
(SSILS)	complementarity	yes	yes
Active-Set Semismooth - Feasibility-enforced		,	,
(ASFLS)	complementarity	ves	yes
Active-Set Semismooth - Feasibility not enforced	,		
(ASILS)	complementarity	yes	yes
Linearly Constrained Lagrangian	pde		
Interior Point Method (ipm)	general	yes	yes

Manual pages for TAO



Topology Abstractions: Distributed Arrays

What is a DM?

Interface for linear algebra to talk to grids
Defines (topological part of) a finite-dimensional function space
 Get an element from this space: DMCreateGlobalVector()
Provides parallel layout
Refinement and coarsening
 DMRefine(), DMCoarsen()
Ghost value coherence
 DMGlobalToLocalBegin()
Matrix preallocation:
 DMCreateMatrix() (formerly DMGetMatrix())

Topology Abstractions

DMDA

Abstracts Cartesian grids in 1, 2, or 3 dimension Supports stencils, communication, reordering Nice for simple finite differences

DMPLEX

Abstracts general topology in any dimension Also supports partitioning, distribution, and global orders Allows aribtrary element shapes and discretizations

DMCOMPOSITE

Composition of two or more DMs

DMNetwork - for discrete networks like power grids and circuits

DMMoab - interface to the MOAB unstructured mesh library

Manual pages for DM

Distributed Array

Interface for topologically structured grids

Defines (topological part of) a finite-dimensional function space

Get an element from this space: DMCreateGlobalVector()

Provides parallel layout

Refinement and coarsening

DMRefineHierarchy()

Ghost value coherence

DMGlobalToLocalBegin()

Matrix preallocation

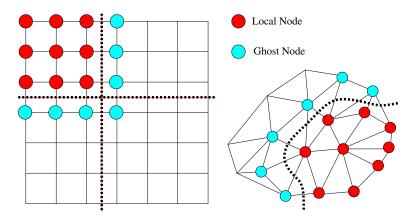
DMCreateMatrix() (formerly DMGetMatrix())

Ghost Values

To evaluate a local function f(x), each process requires

its local portion of the vector x

its ghost values, bordering portions of *x* owned by neighboring processes



DMDA Global Numberings

Proc 2			Proc 3	
25	26	27	28	29
20	21	22	23	24
15	16	17	18	19
10	11	12	13	14
5	6	7	8	9
0	1 2		3	4
Proc 0			Proc 1	
Natural numbering				

Natural numbering

Proc 2			Proc 3		
22	23	28	29		
19	20	26	27		
16	17	24	25		
7	8	13	14		
4	5	11	12		
1	2	9	10		
Proc 0			Proc 1		
	22 19 16 7 4 1	22 23 19 20 16 17 7 8 4 5 1 2	22 23 28 19 20 26 16 17 24 7 8 13 4 5 11 1 2 9		

PETSc numbering

Global: Each vertex has a unique id, belongs on a unique process **Local**: Numbering includes vertices from neighboring processes These are called ghost vertices

Proc 2			Proc 3		
Х	Х	Х	Х	Х	
X	Х	Х	Х	Х	
12	13	14	15	Х	
8	9	10	11	Х	
4	5	6	7	Х	
0	0 1 2		3	Х	
Proc 0			Proc 1		
Local numbering					

Proc 2			Proc 3	
21	22	23	28	29
18	19	20	26	27
15	16	17	24	25
6	7	8	13	14
3	4	5	11	12
0	1	2	9	10
Proc 0			Proc 1	
Global numbering				

DM Vectors

The DM object contains only layout (topology) information

All field data is contained in PETSc ${\tt Vecs}$

Global vectors are parallel

Each process stores a unique local portion

DMCreateGlobalVector(DM dm, Vec *gvec)

Local vectors are sequential (and usually temporary)

Each process stores its local portion plus ghost values DMCreateLocalVector(DM dm, Vec *lvec) includes ghost values!

Coordinate vectors store the mesh geometry

DMDAGetCoordinates (DM dm, Vec *coords) Can be manipulated with their own DMDA DMDAGetCoordinateDA(DM dm,DM *cda)

Updating Ghosts

Two-step Process for Updating Ghosts

enables overlapping computation and communication

DMGlobalToLocalBegin(dm, gvec, mode, lvec)

gvec provides the data mode is either INSERT_VALUES OF ADD_VALUES lvec holds the local and ghost values

DMGlobalToLocalEnd(dm, gvec, mode, lvec)

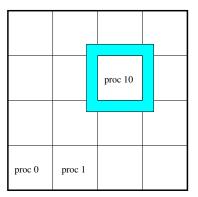
Finishes the communication

Reverse Process

Via DMLocalToGlobalBegin() and DMLocalToGlobalEnd().

DMDA Stencils

Available Stencils



Box Stencil

		proc 10	
proc 0	proc 1		

Star Stencil

Creating a DMDA

xbdy, ybdy: Specifies periodicity or ghost cells

DM_BOUNDARY_NONE, DM_BOUNDARY_GHOSTED, DM_BOUNDARY_MIRROR, DM_BOUNDARY_PERIODIC

type

Specifies stencil: DMDA_STENCIL_BOX or DMDA_STENCIL_STAR

M,N

Number of grid points in x/y-direction

m,n

Number of processes in x/y-direction

dof

Degrees of freedom per node

s

The stencil width

lm,ln

Alternative array of local sizes



Debugging and Profiling

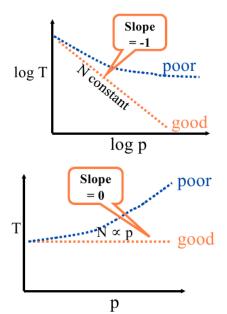
Scalability Definitions

Strong scalability

Fixed problem size execution time T inversely proportional to number of processors p

Weak scalability

Fixed problem size per processor execution time constant as problem size increases



The easiest way to make software scalable is to make it sequentially inefficient. (Gropp 1999)

We really want efficient software

Need a performance model

memory bandwidth and latency algorithmically critical operations (e.g., dot products, scatters) floating point unit

Scalability shows marginal benefit of adding more cores, nothing more

Constants hidden in the choice of algorithm

Constants hidden in implementation

PETSc Debugging

By default, a debug build is provided

Launch the debugger

```
-start_in_debugger [gdb,lldb,dbx,noxterm]
-on_error_attach_debugger [gdb,lldb,dbx,noxterm]
```

Attach the debugger only to some parallel processes

-debugger_nodes 0,1

Set the display (often necessary on a cluster)

-display :0

Debugging Tips

Put a breakpoint in PetscError() to catch errors as they occur

PETSc tracks memory overwrites at both ends of arrays The CHKMEMQ macro causes a check of all allocated memory Track memory overwrites by bracketing them with CHKMEMQ

PETSc checks for leaked memory

Use PetscMalloc() and PetscFree() for all allocation Print unfreed memory on PetscFinalize() with -malloc_dump

Simply the best tool today is Valgrind

It checks memory access, cache performance, memory usage, etc. http://www.valgrind.org Pass -malloc 0 to PETSc when running under Valgrind

 $Might \; need \; - \texttt{trace-children=yes} \; when \; running \; under \; MPI$

--track-origins=yes handy for uninitialized memory

Profiling

Use -log_view [previously -log_summary] for a performance profile
 Event timing
 Event flops
 Memory usage
 MPI messages
Call PetscLogStagePush() and PetscLogStagePop()
 User can add new stages
Call PetscLogEventBegin() and PetscLogEventEnd()
 User can add new events

Call ${\tt PetscLogFlops}$ () to include your flops

Reading -log_view

	Max	Max/Min	Total				
Time (sec):	1.548e+02	1.00122	1.547e+02				
Objects:	1.028e+03	1.00000	1.028e+03				
Flops:	1.519e+10	1.01953	1.505e+10	1.204e+11			
Flops/sec:	9.814e+07	1.01829	9.727e+07	7.782e+08			
MPI Messages:	8.854e+03	1.00556	8.819e+03	7.055e+04			
MPI Message Lengths:	1.936e+08	1.00950	2.185e+04	1.541e+09			
MPI Reductions:	2.799e+03	1.00000					

Also a summary per stage

Memory usage per stage (based on when it was allocated)

Time, messages, reductions, balance, flops per event per stage

Always send -log_view when asking performance questions on mailing list

Event	Count	Time (sec)	Flop	s					GI	.oba	11.			St	ag	e -		Total
	Max Ratio	Max	Ratio	Max	Ratio	Mess	Avg len	Reduct	%T	%F	%M	%L	%R	%T	%F	%M	%L	%R	Mflop/s
Event Stage 1:	Full solve	5																	
VecDot	43 1.0	4.8879e-0	2 8.3 1	L.77e+0	6 1.0	0.0e+00	0.0e+00	4.3e+01	0	0	0	0	0	0	0	0	0	1	73954
VecMDot	1747 1.0	1.3021e+0	0 4.6 8	3.16e+0	7 1.0	0.0e+00	0.0e+00	1.7e+03	0	1	0	0	14	1	1	0	0	27	128346
VecNorm	3972 1.0	1.5460e+0	0 2.5 8	3.48e+0	7 1.0	0.0e+00	0.0e+00	4.0e+03	0	1	0	0	31	1	1	0	0	61	112366
VecScale	3261 1.0	1.6703e-0	1 1.0 3	3.38e+0	7 1.0	0.0e+00	0.0e+00	0.0e+00	0	0	0	0	0	0	0	0	0	0	414021
VecScatterBegin	4503 1.0	4.0440e-0	1 1.0 0	0.00e+0	0.0	6.1e+07	2.0e+03	0.0e+00	0	0	50	26	0	0	0	96	53	0	0
VecScatterEnd	4503 1.0	2.8207e+0	0 6.4 0	0.00e+0	0.0	0.0e+00	0.0e+00	0.0e+00	0	0	0	0	0	0	0	0	0	0	0
MatMult	3001 1.0	3.2634e+0	1 1.1 3	3.68e+0	9 1.1	4.9e+07	2.3e+03	0.0e+00	11	22	40	24	0	22	44	78	49	0	220314
MatMultAdd	604 1.0	6.0195e-0	1 1.0 5	5.66e+0	7 1.0	3.7e+06	1.3e+02	0.0e+00	0	0	3	0	0	0	1	6	0	0	192658
MatMultTranspose	676 1.0	1.3220e+0	0 1.6 6	5.50e+0	7 1.0	4.2e+06	1.4e+02	0.0e+00	0	0	3	0	0	1	1	7	0	0	100638
MatSolve	3020 1.0	2.5957e+0	1 1.0 3	3.25e+0	9 1.0	0.0e+00	0.0e+00	0.0e+00	9	21	0	0	0	18	41	0	0	0	256792
MatCholFctrSym	3 1.0	2.8324e-0	4 1.0 0	0.00e+0	0.0	0.0e+00	0.0e+00	0.0e+00	0	0	0	0	0	0	0	0	0	0	0
MatCholFctrNum	69 1.0	5.7241e+0	0 1.0 6	5.75e+0	8 1.0	0.0e+00	0.0e+00	0.0e+00	2	4	0	0	0	4	9	0	0	0	241671
MatAssemblyBegin	119 1.0	2.8250e+0	0 1.5 0	0.00e+0	0.0	2.1e+06	5.4e+04	3.1e+02	1	0	2	24	2	2	0	3	47	5	0
MatAssemblyEnd	119 1.0	1.9689e+0	0 1.4 0	0.00e+0	0.0	2.8e+05	1.3e+03	6.8e+01	1	0	0	0	1	1	0	0	0	1	0
SNESSolve	4 1.0	1.4302e+0	2 1.0 8	3.11e+0	9 1.0	6.3e+07	3.8e+03	6.3e+03	51	50	52	50	50	991	00	99	100	97	113626
SNESLineSearch	43 1.0	1.5116e+0	1 1.0 1	L.05e+0	8 1.1	2.4e+06	3.6e+03	1.8e+02	5	1	2	2	1	10	1	4	4	3	13592
SNESFunctionEval	55 1.0	1.4930e+0	1 1.0 0	0.00e+0	0.0	1.8e+06	3.3e+03	8.0e+00	5	0	1	1	0	10	0	3	3	0	0
SNESJacobianEval	43 1.0	3.7077e+0	1 1.0 7	7.77e+0	6 1.0	4.3e+06	2.6e+04	3.0e+02	13	0	4	24	2	26	0	7	48	5	429
KSPGMRESOrthog	1747 1.0	1.5737e+0	0 2.9 1	L.63e+0	8 1.0	0.0e+00	0.0e+00	1.7e+03	1	1	0	0	14	1	2	0	0	27	212399
KSPSetup	224 1.0	2.1040e-0	2 1.0 0	0.00e+0	0.0	0.0e+00	0.0e+00	3.0e+01	0	0	0	0	0	0	0	0	0	0	0
KSPSolve	43 1.0	8.9988e+0	1 1.0 7	7.99e+0	9 1.0	5.6e+07	2.0e+03	5.8e+03	32	49	46	24	46	62	99	88	48	88	178078
PCSetUp	112 1.0	1.7354e+0	1 1.0 6	5.75e+0	8 1.0	0.0e+00	0.0e+00	8.7e+01	6	4	0	0	1	12	9	0	0	1	79715
PCSetUpOnBlocks	1208 1.0	5.8182e+0	0 1.0 6	5.75e+0	8 1.0	0.0e+00	0.0e+00	8.7e+01	2	4	0	0	1	4	9	0	0	1	237761
PCApply	276 1.0	7.1497e+0	1 1.0 7	7.14e+0	9 1.0	5.2e+07	1.8e+03	5.1e+03	25	44	42	20	41	49	88	81	39	79	200691

Communication Costs

Reductions: usually part of Krylov method, latency limited

VecDot

VecMDot

VecNorm

MatAssemblyBegin

Change algorithm (e.g. IBCGS)

Point-to-point (nearest neighbor), latency or bandwidth

VecScatter MatMult PCApply MatAssembly SNESFunctionEval SNESJacobianEval Compute subdomain boundary fluxes redundantly Ghost exchange for all fields at once Better partition

Conclusions

PETSc can help You

Solve algebraic and DAE problems in your application area Rapidly develop efficient parallel code, can start from examples Develop new solution methods and data structures Debug and analyze performance Advice on software design, solution algorithms, and performance $petsc-{users, dev, maint}@mcs.anl.gov}$

You can help PETSc

Report bugs and inconsistencies, or if you think there is a better way Tell us if the documentation is inconsistent or unclear Consider developing new algebraic methods as plugins, contribute if your idea works

Follow Up; Getting Help

http://www.mcs.anl.gov/petsc

Public questions: petsc-users@mcs.anl.gov, with searchable archives Private questions: petsc-maint@mcs.anl.gov

Hands On Exercises

Instructions http://www.mcs.anl.gov/petsc/petsc-3.7-atpesc2016/tutorials/HandsOnExercise.html Examples Linear Poisson equation on a 2D grid src/ksp/ksp/examples/tutorials/ex50.c Nonlinear ODE arising from a time-dependent one-dimensional PDE src/ts/examples/tutorials/ex2.c Nonlinear PDE on a structured grid src/snes/examples/tutorials/ex19.c Linear Stokes-type PDE on a structured grid src/ksp/ksp/examples/tutorials/ex42.c Nonlinear time-dependent PDE on Unstructured Grid src/ts/examples/tutorials/ex11.c